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# C++ Programme for total dominator chromatic number of ladder graphs through simple transformations

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#### Abstract

A total dominator coloring of a graph  $\mathbb{G} = (\mathbb{V}, \mathbb{E})$  without isolated vertices, along with each vertex in  $\mathbb{G}$ , is a proper coloring that dominates a color class. The total chromatic dominator number of  $\mathbb{G}$  is the minimum number of color classes with further assumption that each vertex in  $\mathbb{G}$  dominates a color class properly and is represented as  $\chi_{td}(\mathbb{G})$ . In this manuscript, we consider the chromatic total dominator number of ladder graphs through fundamental transformations via the program C++.

#### **Keywords**

Coloring, Total dominator coloring, Total dominator chromatic number.

AMS Subject Classification 05C69, 68W25.

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## 1. Introduction

We mainly find ladder graphs in this manuscript. For additional information in graph theory and its applications, we suggest the reader to refer F. Harrary [4]. Allow  $\mathbb{G} = (\mathbb{V}, \mathbb{E})$ be a graph without isolated vertices. For any two graphs  $\mathbb{G}$  and  $\mathbb{H}$ , we characterize the cartesian product, signified by  $\mathbb{G} \times \mathbb{H}$ , to be the graph with vertex set  $\mathbb{V}(\mathbb{G}) \times \mathbb{V}(\mathbb{H})$  and edges between two vertices  $(u_1, v_1)$  and  $(u_2, v_2)$  iff either  $u_1 = u_2$ and  $v_1v_2 \in E(H)$  or  $u_1u_2 \in E(G)$  and  $v_1 = v_2$ .

In general, for  $n \ge 2$ , we characterize a ladder graph as  $P_2 \times P_n$  and is signified by  $L_n$  and  $|\mathbb{V}(L_n)| = p = 2n, n \ge 2$ .

A proper coloring of  $\mathbb G$  is an assignment of colors to the vertices of  $\mathbb G,$  in a way that adjacent vertices have different

colors. The smallest number of colors for which  $\mathbb{G}$  is properly colored is considered a chromatic number of  $\mathbb{G}$ , and  $\chi(\mathbb{G})$  is denoted. A total dominator coloring (*td*-coloring) of  $\mathbb{G}$  is a proper coloring of  $\mathbb{G}$  with additional axioms that is properly dominated color class by every vertex in  $\mathbb{G}$ . Let  $\chi_{td}(\mathbb{G})$  be the total dominator chromatic number and is defined by the minimum number of colors needed in a total dominator coloring of  $\mathbb{G}$ . This principle was developed in [1] by Vijayalekshmi. This thought is often pointed to as a  $\mathbb{G}$ ,  $(k \ge 1)$ Smarandachely k-dominator color and was presented in [2] by Vijayalekshmi. A Smarandachely *k*-dominator coloring of  $\mathbb{G}$  for an integer  $k \ge 1$  is a proper coloring of G, so that each vertex in a  $\mathbb{G}$  graph properly dominates a color class of k. The smallest number of colors for which there exists a Smarandachely k-dominator coloring of  $\mathbb{G}$  is called the Smarandachely k-dominator chromatic number of  $\mathbb{G}$  and is denoted by  $\chi_{td}^s(\mathbb{G})$ .

Let  $\mathscr{C}$  be a minimum *td*-coloring of  $\mathbb{G}$ . We say a color class is considered a non-dominated color class (n - d color class) if no vertex of  $\mathbb{G}$  dominates it and these color classes are often considered repeated color classes.

We recommend the author to pertain to [3, 5, 6] for further information on this theory and its applications.

## 2. Preliminaries

In this segment, we remember the critical [3] theorem which is quite helpful in our research. For the subsequent observation the minimum dominator chromatic number of ladder graphs has been identified.

For every  $n \ge 2$ , the total dominator chromatic number of a ladder graph is

$$\chi_{td} (\mathscr{C}_n) = \begin{cases} 2\lfloor \frac{p}{6} \rfloor + 2, & ifn \equiv 0 (mod6) \\ 2\lfloor \frac{p-2}{6} \rfloor + 4, & \\ 2\lfloor \frac{p-4}{6} \rfloor + 4, & otherwise. \end{cases}$$

In this manuscript we obtain a C++ program which uses fundamental transformations to find the *td*-chromatic number of ladder graphs.

## 3. Main Result

In this section, We have to find the total dominator chromatic number of ladder graphs using C++ programme. The C++ programme is successfully compiled and run on C++ platform. The runtime test is included.

#### Programme as follows

```
#include "stdafx.h"
#include <Windows.h>
#include <conio.h>
#include <iostream>
using namespace std;
int main() {
int inpt;
cout << "Enter the Value of Ln" << endl;</pre>
cin >> inpt;
int N = inpt + inpt; int M = inpt + inpt;
int** ary = new int*[N];
int** mat = new int*[N];
int** mat1 = new int*[N];
int** mat2 = new int*[N];
int** mat3 = new int*[N];
int** matsum = new int*[N];
for (int i = 0; i < N; ++i)
{
  ary[i] = new int[M]; mat[i] = new int[M]; mat1[i] = new int[M];
mat2[i] = new int[M]; mat3[i] = new int[M]; matsum[i] = new int[M];
}
int k, l, sum;
HANDLE p = GetStdHandle(STD_OUTPUT_HANDLE);
SetConsoleTextAttribute(p, FOREGROUND INTENSITY | FOREGROUND INTENSITY);
for (int i = 0; i < N; ++i)
for (int j = 0; j < M; ++j)
ary[i][j] = i;
cout << "\n" << "The Adjacency Matrix for L" << inpt << "\n";
for (int i = 0; i < N; i++)
{
if (i % 2 == 0)
{
for (int j = 0; j < N; j++)
{
if (ary[j][i] == i + 1 | ary[j][i] == i - 1 | ary[j][i] == i + 3)
{
mat[i][j] = 1;
cout << mat[i][j] << " ";</pre>
}
```



```
else
{
mat[i][j] = 0;
cout << mat[i][j] << " ";</pre>
}
}
}
else
{
for (int j = 0; j < N; j++)</pre>
{
if (ary[j][i] == i + 1 | ary[j][i] == i - 1 | ary[j][i] == i - 3)
{
mat[i][j] = 1;
cout << mat[i][j] << " ";</pre>
}
else
{
mat[i][j] = 0;
cout << mat[i][j] << " ";</pre>
}
}
}
cout << "\n";</pre>
}
cout << "\n" << "ADJACENCY MATRIX BY SUBSTRATING THE ROW ASSENDING VALUES" << "\n";
for (int i = 0; i < N; i++)
{
int sum = 0;
for (int j = 0; j < N; j++)
{
if (i >= 2 && i <= 5 && mat[i][j] == 1 && mat1[i - 2][j] == 1)
{
mat1[i][j] = mat[i][j] - mat[i - 2][j];
}
else if (i >= 6 && mat[i][j] == 1 && mat1[i - 2][j] == 1 && matsum[i - 4][0] != 1)
{
mat1[i][j] = mat[i][j] - mat[i][j];
}
else if (i >= 4 && mat[i][j] == 1 && mat1[i - 2][j] == 0 && matsum[i - 4][0] == 1)
{
mat1[i][j] = mat[i][j] - mat1[i - 4][j];
}
else
{
mat1[i][j] = mat[i][j];
}
sum = sum + mat1[i][j];
}
matsum[i][0] = sum;
}
for (int i = 0; i < N; i++)
{
for (int j = 0; j < N; j++)
{
if (mat1[i][j] == 1)
```



```
{
SetConsoleTextAttribute(p, FOREGROUND RED | FOREGROUND INTENSITY);
cout << mat1[i][j] << " ";</pre>
}
else
{
SetConsoleTextAttribute(p, FOREGROUND_INTENSITY | FOREGROUND_INTENSITY);
cout << mat1[i][j] << " ";
}
}
cout << "\n";</pre>
}
SetConsoleTextAttribute(p, FOREGROUND_INTENSITY | FOREGROUND_INTENSITY);
cout << "\n"<< "ADJACENCY MATRIX BY SUBSTRATING THE COLUMN VALUES";
if (N%3 == 0)
{
N = N - 1;
for (int i = N; i \ge 0; i--)
{
for (int j = N; j \ge 0; j--)
{
if (i >= 4 && mat1[i][j] == 1 && mat1[i - 4][j] == 1)
{
mat1[i - 4][j] = mat1[i][j] - mat1[i - 4][j];
}
else if (i >= 2 && mat1[i][j] == 1 && mat1[i - 2][j] == 1)
{
mat1[i - 2][j] = mat1[i][j] - mat1[i - 2][j];
}
}
}
N = N + 1;
cout << "\n";</pre>
}
else
{
for (int i = 0; i < N; i++)
{
for (int j = 0; j < N; j++)
{
if (mat1[i][j] == 1 && mat1[i][j+2] == 1)
{
mat1[i][j+2] = mat1[i][j+2] - mat1[i][j];
}
else if (mat1[i][j] == 1 && mat1[i][j + 4] == 1)
{
mat1[i][j + 4] = mat1[i][j + 4] - mat1[i][j];
}
else
{
mat1[i][j] = mat1[i][j];
}
}
}
cout << "\n";
}
```



```
for (int i = 0; i < N; i++)
{
for (int j = 0; j < N; j++)
{
if (mat1[i][j] == 1)
{
SetConsoleTextAttribute(p, FOREGROUND_RED | FOREGROUND_INTENSITY);
cout << mat1[i][j] << " ";</pre>
}
else
{
SetConsoleTextAttribute(p, FOREGROUND_INTENSITY | FOREGROUND_INTENSITY);
cout << mat1[i][j] << " ";
}
}
cout << "\n";</pre>
}
SetConsoleTextAttribute(p, FOREGROUND INTENSITY | FOREGROUND INTENSITY);
cout << "\n";</pre>
int ary2[] = \{ 0, 1, 4, 5, 2, 3 \}, aaa = 0;
for (int i = 0; i < N; i++)
{
for (int j = 0; j < N; j++)
{
if (ary2[aaa] > N-1)
{
ary2[aaa] = ary2[aaa]-2;
mat3[i][j] = mat1[ary2[aaa]][j];
}
mat3[i][j] = mat1[ary2[aaa]][j];
}
if (aaa < 5)
{
ary2[aaa] = ary2[aaa] + 6;
aaa = aaa + 1;
}
else if (aaa = 5)
{
ary2[aaa] = ary2[aaa] + 6;
aaa = 0;
}
}
cout << "FINAL SUB MATRIXES AFTER SUBSTRACTING THE COLUMN FROM BOTTOM TO TOP
" << "\n";
k = 0;
for (int i = 0; i < N; i++)
{
for (int j = 0; j < N; j++)
{
if (j % 2 == 0 && i % 2 == 0 && mat1[i][j] == 0 && mat1[i][j + 1] == 1 ||
  mat1[i][j] == 1)
{
SetConsoleTextAttribute(p, FOREGROUND_RED | FOREGROUND_INTENSITY);
cout << mat1[i][j] << " ";
}
else if (j % 2 != 0 && i % 2 != 0 && mat1[i][j] == 0 && mat1[i][j - 1] == 1)
```



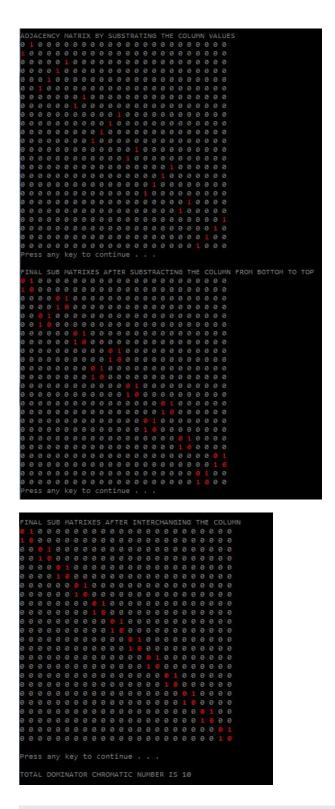
```
{
SetConsoleTextAttribute(p, FOREGROUND RED | FOREGROUND INTENSITY);
k = k + 1;
cout << mat1[i][j] << " ";</pre>
}
else
{
SetConsoleTextAttribute(p, FOREGROUND INTENSITY | FOREGROUND INTENSITY);
cout << mat1[i][j] << " ";</pre>
}
}
cout << "\n";</pre>
}
SetConsoleTextAttribute(p, FOREGROUND_INTENSITY | FOREGROUND_INTENSITY);
cout << "\n" << "FINAL SUB MATRIXES AFTER INTERCHANGING THE COLUMN" << "\n";
for (int i = 0; i < N; i++)
{
for (int j = 0; j < N; j++)
{
if ( i% 2==0 && mat3[i][j] == 0 && mat3[i][j + 1] == 1 || mat3[i][j] == 1)
{
SetConsoleTextAttribute(p, FOREGROUND_RED | FOREGROUND_INTENSITY);
cout << mat3[i][j] << " ";</pre>
}
else if (i % 2 != 0 && mat3[i][j] == 0 && mat3[i][j - 1] == 1)
{
SetConsoleTextAttribute(p, FOREGROUND_RED | FOREGROUND_INTENSITY);
cout << mat3[i][j] << " ";</pre>
}
else
{
SetConsoleTextAttribute(p, FOREGROUND_INTENSITY | FOREGROUND_INTENSITY);
cout << mat3[i][j] << " ";</pre>
}
}
cout << "\n";
SetConsoleTextAttribute(p, FOREGROUND_INTENSITY | FOREGROUND_INTENSITY);
}
cout << "\n";</pre>
if (inpt % 3 == 0)
{
cout <<"\n" << "TOTAL DOMINATOR CHROMATIC NUMBER IS " << (2 * (k / 3)) + 2
<< "\n";
}
else
{
cout << "\n" << "TOTAL DOMINATOR CHROMATIC NUMBER IS " << (2 * (k-1)/3) + 4
<< "\n";
}
system("Pause");
return 0;
for (int i = 0; i < N; ++i)
delete[] ary[i], ary, mat1[i], mat1, mat[i], mat, matsum[i], matsum, mat2[i], mat2,
mat3[i], mat3;
}
```



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return 0; }

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	cy Matrix for L12
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	3



# 4. Conclusion

Within this manuscript, we treat the total dominator chro-



matic number of ladder graphs in a simplified and enhanced fashion utilizing elementary transformations by C++ programme.

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